

# Specification, Semantics, and Verification of Quantum Programs

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University of Vermont

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Certified Programs and Proofs (CPP)  
Jan 13, 2026



```
# input: s is a bitvector given as a string
def bernstein_vazarani(s):
    n = len(s)

    qiskit_circuit = QuantumCircuit(n + 1)

    qiskit_circuit.x(n)
    qiskit_circuit.barrier()
    qiskit_circuit.h(range(n + 1))
```

```
for ii, yesno in enumerate(reversed(s)):
    if yesno == "1":
        qiskit_circuit.cx(ii, n)

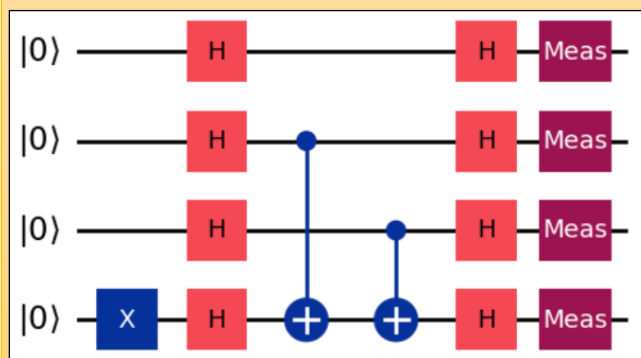
qiskit_circuit.barrier()
qiskit_circuit.h(range(n + 1))

qiskit_circuit.barrier()
for i in range(n+1):
    qiskit_circuit.append(measure, [i])

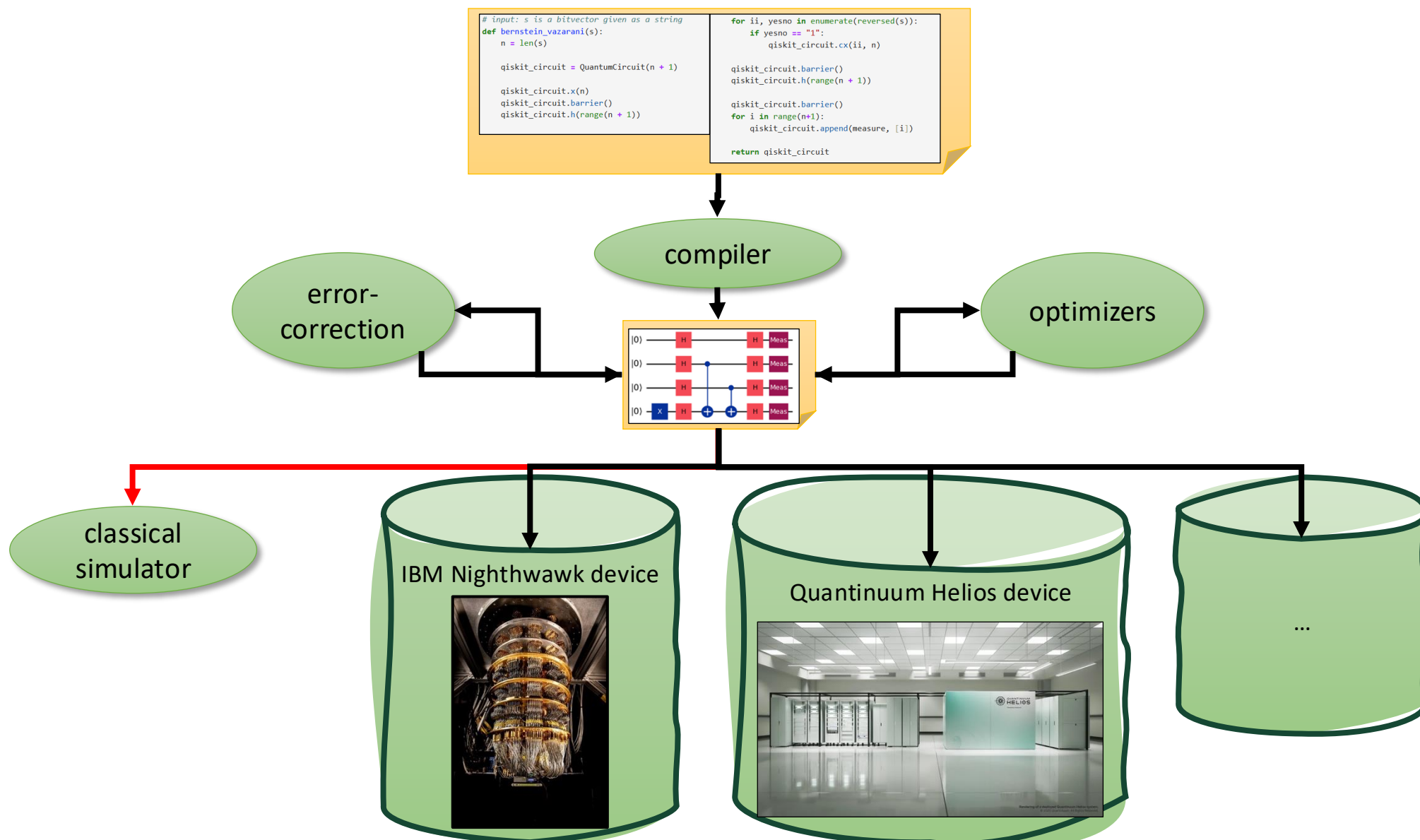
return qiskit_circuit
```

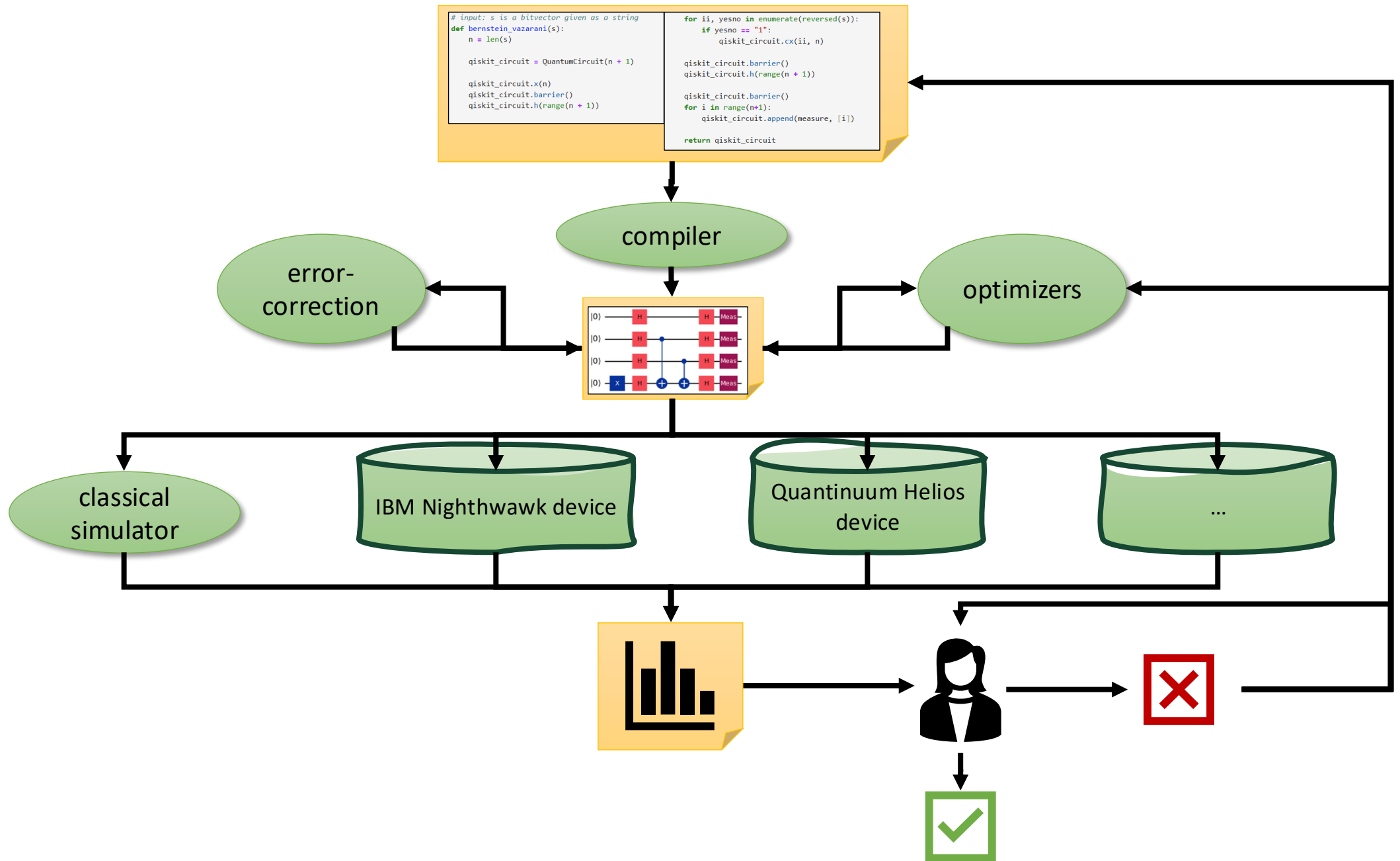
quantum  
DSL

compiler



quantum circuit





# Outline

0. Vocabulary
1. Semantics of quantum circuits
2. Case study: specifying and verifying quantum programs in Rocq
3. Case study: verifying error correction and fault tolerance
4. Wrap up

# Vocab Lesson



Formal Verification

machine-aided  
mathematical & logical  
methods to prove a  
program matches its  
specification

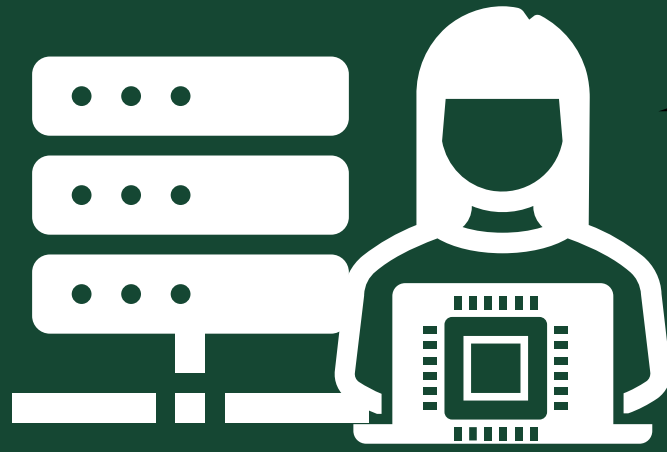
# Vocab Lesson



Quantum Verification

does a quantum experiment demonstrate behavior that can't be achieved by a classical system?

# Vocab Lesson

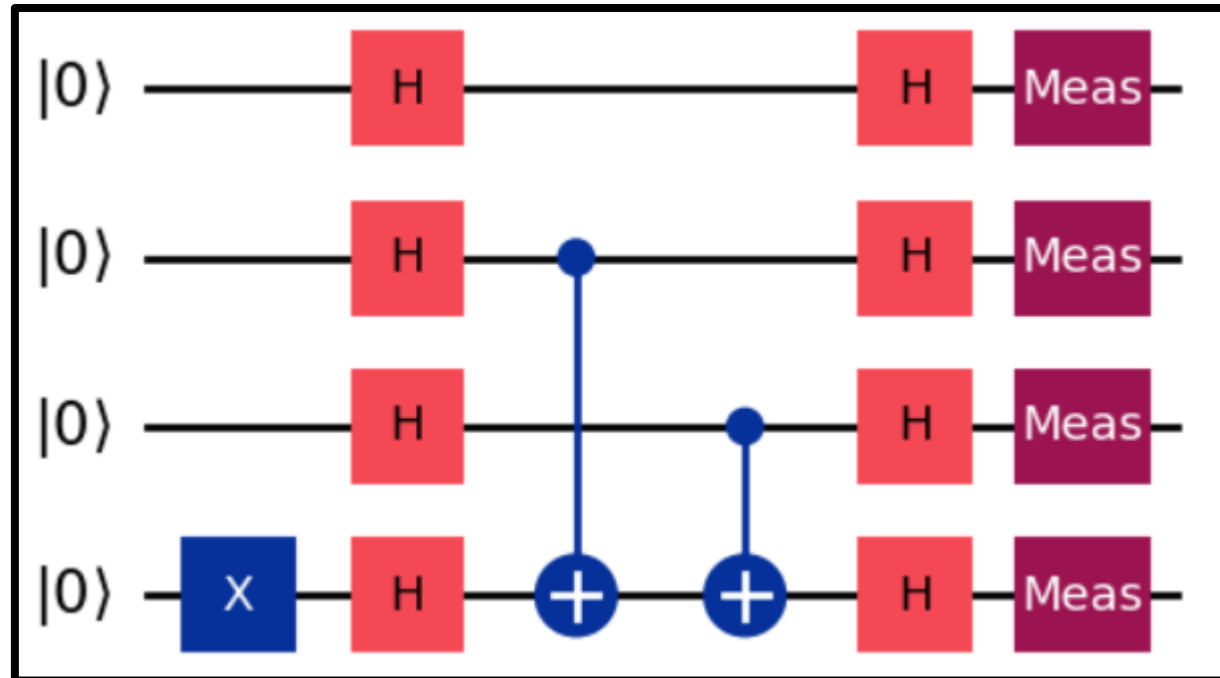


Verification & Validation

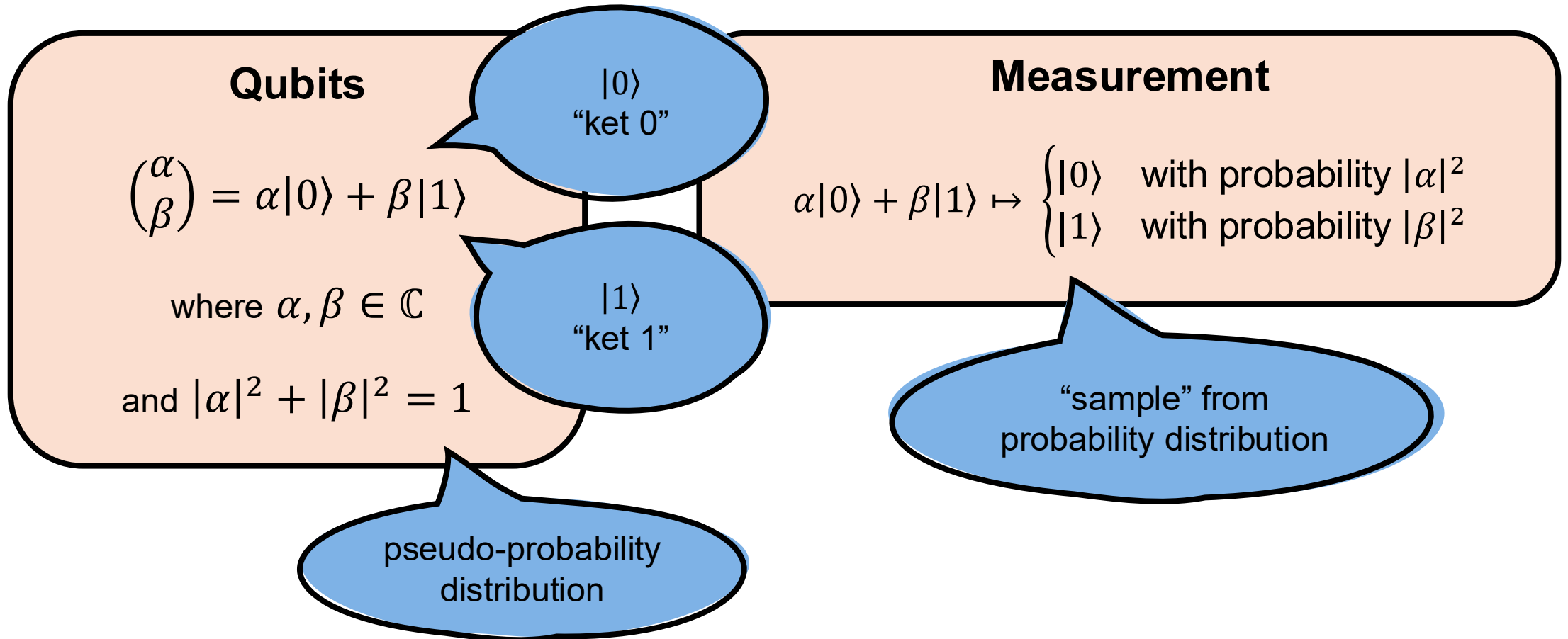
verification  
=  
model checking + formal  
verification,  
OR  
simulation, timing  
analysis, testing

# Semantics of a quantum circuit

# Quantum circuits



# Quantum circuits



# Quantum circuits

## Multi-qubit systems

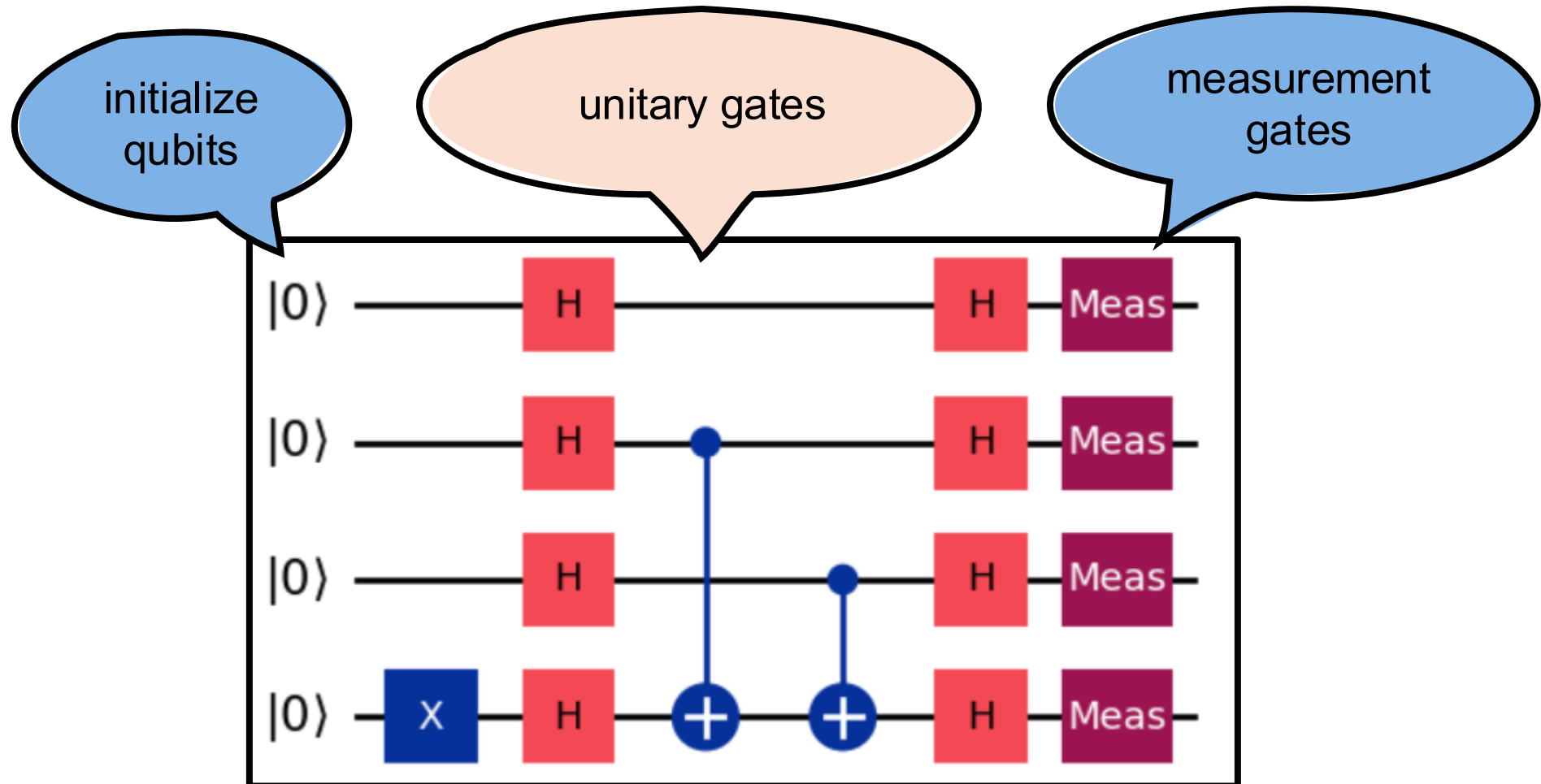
$$\alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle$$

where  $\alpha_{ij} \in \mathbb{C}$

$$\text{and } \sum |\alpha_{ij}|^2 = 1$$

$n$  qubits  
=  
complex vector of size  $2^n$

# Quantum circuits



# Quantum circuits

## Unitary transformations

Square complex matrix  $U \in \mathbb{C}^{2^n \times 2^n}$

such that  $U^{-1} = U^\dagger$

# Quantum circuits

## Qubits

$$\begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \alpha|0\rangle + \beta|1\rangle$$

where  $\alpha, \beta \in \mathbb{C}$

$$\text{and } |\alpha|^2 + |\beta|^2 = 1$$

## Initialization

$$|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \quad |1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

## Measurement

$$\alpha|0\rangle + \beta|1\rangle \mapsto \begin{cases} |0\rangle & \text{with probability } |\alpha|^2 \\ |1\rangle & \text{with probability } |\beta|^2 \end{cases}$$

## Unitary matrices

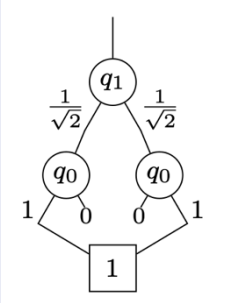
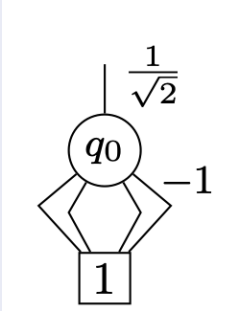
Square complex matrix  $U \in \mathbb{C}^{2^n \times 2^n}$

such that  $U^{-1} = U^\dagger$

# Open Problem #1: Representations of circuit semantics

Representation	$ +\rangle$ state	Hadamard gate
vectors and matrices	$\begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{pmatrix}$	$\begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix}$
Hilbert spaces and linear transformations	$\frac{1}{\sqrt{2}} 0\rangle + \frac{1}{\sqrt{2}} 1\rangle$	$ 0\rangle \mapsto \frac{1}{\sqrt{2}} 0\rangle + \frac{1}{\sqrt{2}} 1\rangle$ $ 1\rangle \mapsto \frac{1}{\sqrt{2}} 0\rangle - \frac{1}{\sqrt{2}} 1\rangle$

# Open Problem #1: Representations of circuit semantics

Representation	+⟩ state	Hadamard gate
vectors and matrices	$\begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{pmatrix}$	$\begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix}$
Quantum Decision Diagrams (QDDs)		

D. M. Miller and M. A. Thornton.

**QMDD: A Decision Diagram Structure for Reversible and Quantum Circuits.**

ISMVL 2006

R. Wille, S. Hillmich, and L. Burgholzer.

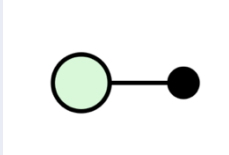
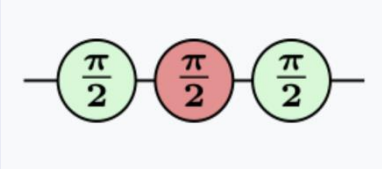
**Decision Diagrams for Quantum Computing.**

In Design Automation of Quantum Computers, 2023.

# Open Problem #1: Representations of circuit semantics

Representation	$ +\rangle$ state	Hadamard gate
vectors and matrices	$\begin{pmatrix} \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} \end{pmatrix}$	$\begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix}$
Phase polynomials	$\frac{1}{\sqrt{2}} \sum_{x \in \mathbb{Z}_2}  x\rangle$	$ x\rangle \mapsto \frac{1}{\sqrt{2}} \sum_{y \in \mathbb{Z}_2} e^{2\pi i \frac{xy}{2}}  y\rangle$

# Open Problem #1: Representations of circuit semantics


Representation	$ +\rangle$ state	Hadamard gate
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ZX diagrams		

# Case Study: Specifying and verifying quantum programs in Rocq




## INQWIRE

Verified Software for the Computers of Tomorrow

-  **QWIRE** Public ⋮


A quantum circuit language and formal verification tool

● Coq ☆ 106 👤 27

 **SQIR** Public ⋮


A Small Quantum Intermediate Representation

● Rocq Prover ☆ 91 👤 24

 **QuantumLib** Public ⋮

Coq library for reasoning about quantum programs

● Rocq Prover ☆ 40 👤 12

 **VyZX** Public ⋮

Verifying the ZX Calculus

● Rocq Prover ☆ 20 👤 4



Robert Rand  
University of Chicago



### QWIRE: a core language for quantum circuits.

Jennifer Paykin, Robert Rand, and Steve Zdancewic.  
POPL 2017

### QWIRE practice: formal verification of quantum circuits in Coq.

Robert Rand, Jennifer Paykin, and Steve Zdancewic.  
QPL 2017.

### ReQWIRE: Reasoning about Reversible Quantum Circuits.

Robert Rand, Jennifer Paykin, Dong-Ho Lee, and Steve Zdancewic.  
QPL 2018.

### Proving Quantum Programs Correct.

Kesha Hietala, Robert Rand, Shih-Han Hung, Liyi Li, Michael Hicks.  
ITP 2021

### A Verified Optimizer for Quantum Circuits.

Kesha Hietala, Robert Rand, Shih-Han Hung, Xiaodi Wu, Michael Hicks.  
POPL 2021.

### QuantumLib: A Library for Quantum Computing in Coq.

Jacob Zweifler, Kesha Hietala, Robert Rand  
Coq Workshop 2022.

### A Formally Certified End-to-End Implementation of Shor's Factorization Algorithm.

Yuxiang Peng, Kesha Hietala, Runzhou Tao, Liyi Li, Robert Rand, Michael Hicks, Xiaodi Wu. PNAS 2023

### VyZX: Formal Verification of a Graphical Quantum Language.

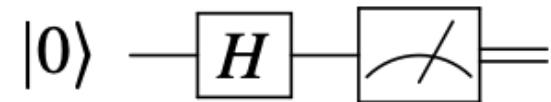
Adrian Lehmann, Ben Caldwell, Bhakti Shah, Robert Rand.  
2023.

# QWIRE

quantum circuit DSL

```
Definition coin_flip : Box One Bit := box_ () => meas $ _H $ init0 $ ().
```

$\begin{cases} |0\rangle & \text{with prob. 0.5} \\ |1\rangle & \text{with prob. 0.5} \end{cases}$



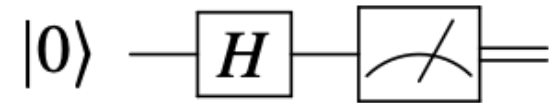
# Handling measurement probabilities

Qubits	Semantics	Measurement
pure state $ \varphi\rangle \in \mathbb{C}^{2^n}$	$U \in \mathbb{C}^{2^n \times 2^n}$	probabilistic step $ \varphi\rangle \rightarrow^p  \varphi'\rangle$
mixed state <i>(probability distribution over pure states)</i> $\rho = \sum_i p_i  \varphi_i\rangle\langle\varphi_i  \in \mathbb{C}^{2^n \times 2^n}$	completely positive trace-preserving (CPTP) maps $\rho \mapsto \rho'$	deterministic step $\rho \rightarrow \rho'$

# QWIRE

```
Definition coin_flip : Box One Bit :=  
  box_ () => meas $ _H $ init0 $ ().
```

$\begin{cases} |0\rangle & \text{with prob. } 0.5 \\ |1\rangle & \text{with prob. } 0.5 \end{cases}$



As a mixed state:  
 $0.5|0\rangle\langle 0| + 0.5|1\rangle\langle 1|$

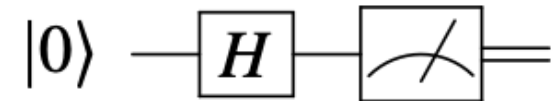
# QWIRE

```
Definition coin_flip : Box One Bit :=  
  box_ () => meas $ _H $ init0 $ ().
```

```
Lemma fair_toss : [[coin_flip]] (I 1) = 0.5 * |0><0| + 0.5 * |1><1|.   
Proof.  
  matrix_denote. Msimpl. solve_matrix.  
Qed.
```

custom tactics

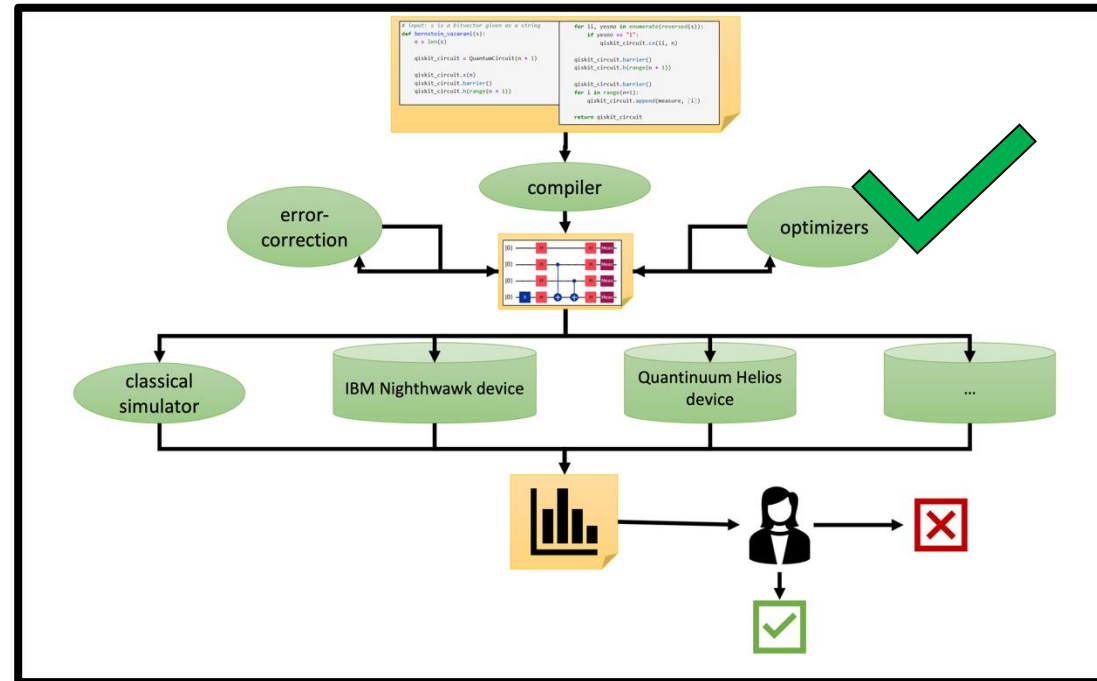
$\begin{cases} |0\rangle & \text{with prob. } 0.5 \\ |1\rangle & \text{with prob. } 0.5 \end{cases}$



As a mixed state:

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# VOQC



## A Verified Optimizer for Quantum Circuits.

Kesha Hietala, Robert Rand, Shih-Han Hung, Xiaodi Wu, Michael Hicks.  
POPL 2021.

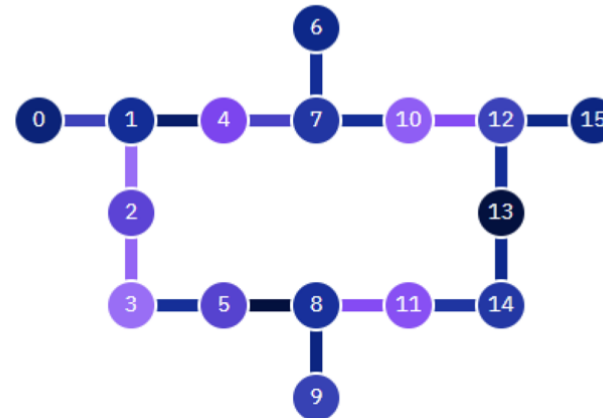
# VOQC

- Verified optimizations

**Lemma** `zyz_to_zyz_correct` :  $\forall \theta_1 \xi \theta_2 \xi_1 \theta \xi_2,$   
 $\text{zyz\_to\_zyz } \theta_1 \xi \theta_2 = (\xi_1, \theta, \xi_2) \rightarrow$   
 $\text{y\_rotation } \theta_2 \times \text{phase\_shift } \xi \times \text{y\_rotation } \theta_1$   
 $\propto \text{phase\_shift } \xi_2 \times \text{y\_rotation } \theta \times \text{phase\_shift } \xi_1.$

- Verified hardware-aware passes

- circuit mapping
- layout
- routing



# Open Problem #2: programs vs circuits

```
# input: s is a bitvector given as a string
def bernstein_vazarani(s):
    n = len(s)

    qiskit_circuit = QuantumCircuit(n + 1)

    qiskit_circuit.x(n)
    qiskit_circuit.barrier()
    qiskit_circuit.h(range(n + 1))

    for ii, yesno in enumerate(reversed(s)):
        if yesno == "1":
            qiskit_circuit.cx(ii, n)

    qiskit_circuit.barrier()
    qiskit_circuit.h(range(n + 1))

    for i in range(n+1):
        qiskit_circuit.append(measure, [i])

    return qiskit_circuit
```

loops/classical control flow

parameterized by number of qubits, other inputs

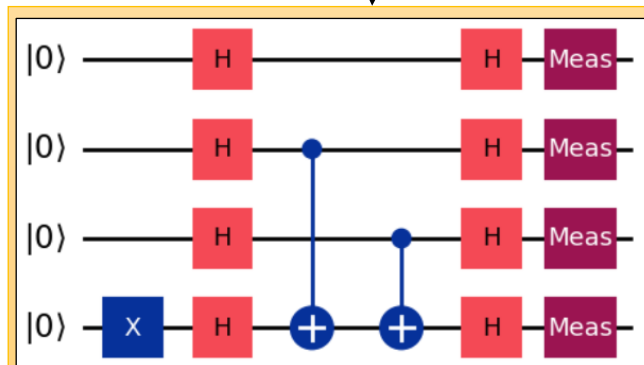
quantum DSL

higher-order quantum abstractions

interleaved classical-quantum computation

compiler

Proto-Quipper family



quantum circuit

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loops/classical control flow

parameterized by number of qubits, other inputs

quantum DSL

higher-order quantum abstractions

interleaved classical-quantum computation

## Verifying quantum programs @ POPL week

Jyun-Ao Lin et al.  
**Verifying Repeat-Until-Success Circuits with AutoQ.**  
PLanQC 2026.

Yingte Xu.  
**A Unified Assertion-Based Framework for Classical-Quantum Program Verification**  
PLanQC 2026.

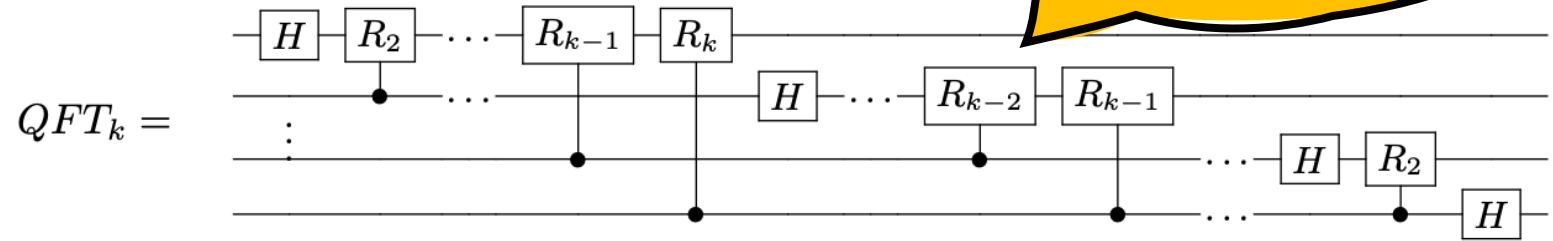
Robert Booth and Cole Comfort.  
**Denotational Semantics for Stabiliser Quantum Programs.**  
PLanQC 2026.

Parosh Aziz Abdulla et al.  
**Parameterized Verification of Quantum Circuits.**  
POPL 2026.

Ken Sakayori, Andrea Colledan, and Ugo Dal Lago.  
**On Circuit Description Languages, Indexed Monads, and Resource Analysis.**  
POPL 2026.

# SQIR – Simple Quantum IR

True for any size circuit and any basis state



```

Fixpoint controlled_rotations n :=
match n with
| 0 | 1 => SKIP
| S n' => controlled_rotations n' ; control n' (Rz (2π / 2n) 0)
end.
Fixpoint QFT n :
match n with
| 0 => SKIP
| S n' => H 0 ; controlled_rotations n ; map_qubits (fun q => q + 1) (QFT n')
end.
    
```

**Lemma**  $QFT\_semantics : \forall n\ x, n > 0 \rightarrow \llbracket QFT\ n \rrbracket_n \times |x\rangle = \frac{1}{\sqrt{2^n}} \bigotimes_{j=0}^{n-1} (|0\rangle + e^{2\pi i x / 2^{n-j}} |1\rangle).$

As a unitary on n qubits...

...QFT takes the basis state  $|x\rangle$ ...

...to the following pure state.

# Tradeoffs

Pros	Cons
Powerful reasoning techniques	Requires expert proof developers, limited automation

## Quantum Automatic Verification @ POPL Week

Xiaoguan Xu, Li Zhou, Mingsheng Ying.

**VC-Qiskit: Automated–Interactive Verification of Qiskit Passes with Minimal Intrusion**

Christophe Charetton.

**Democratizing quantum formal verification: the path-sum way**

Robin Adams, Jean-Philippe Bernardy, Lorenzo Perticonne, and Jeremy Pope.

**A Graded Modal Type Theory for Pulse Schedules**

Xingte Xu.

**A Unified Assertion-Based Framework for Classical-Quantum Program Verification**

Kayo Tei, Haruto Minshina, Naoki Yamamoto, and Kazunori Ueda.

**Graph Rewriting Language as a Platform for Quantum Diagrammatic Calculi**

Marco Lewis and Benoit Valiron

**Finding Photonics Circuits via  $\delta$ -weakening SMT**

Parosh Aziz Abdulla et al.

**Parameterized Verification of Quantum Circuits**



Open Problem #3

# Tradeoffs

Pros	Cons
Powerful reasoning techniques	Requires expert proof developers, limited automation
Statement of equivalence conceptually clear	Other quantum specifications unintuitive

Open Problem #4

Christophe Chareton

**Democratizing quantum  
formal verification: the  
path-sum way.**

PLanQC keynote

# Tradeoffs

Pros	Cons
Powerful reasoning techniques	Requires expert proof developers, limited automation
Statement of equivalence conceptually clear	Other quantum specifications unintuitive
Can reason about complex properties e.g. ancillae and oracles	Not straightforward to reason about errors, probabilistic algorithms

## Quantum Errors @ POPL Week

Anurudh Peduri, Jam Kabeer, Ali Khan, Gilles Barthe, Michael Walter

**Traq: Estimating the Quantum Cost of Classical Programs**

Robin Adams, Jean-Philippe Bernardy, Lorenzo Perticonne, and Jeremy Pope.

**A Graded Modal Type Theory for Pulse Schedules**

Jane Moore, Michael Hart, and John McAllister.

**Efficient Parallel Compilation and Profiling of Quantum Circuits at Large Scales**

Yu-Hsuan Wu, Yue Shi, Junyi Liu, and Yuxiang Peng.

**A Pulse-Level DSL for Real-Time Quantum Control with Hardware Compilation and Emulation**



Open Problem #5

# Case Study: Verifying Error Correction and Fault Tolerance

## Verifying Fault-Tolerance of Quantum Error Correction Codes

Check for  
update

Kean Chen<sup>1</sup>✉<sup>ID</sup>, Yuhao Liu<sup>1</sup><sup>ID</sup>, Wang Fang<sup>2</sup><sup>ID</sup>, Jennifer Paykin<sup>3</sup>,  
Xin-Chuan Wu<sup>4</sup>, Albert Schmitz<sup>3</sup>, Steve Zdancewic<sup>1</sup>✉<sup>ID</sup>, and Gushu Li<sup>1</sup>✉<sup>ID</sup>



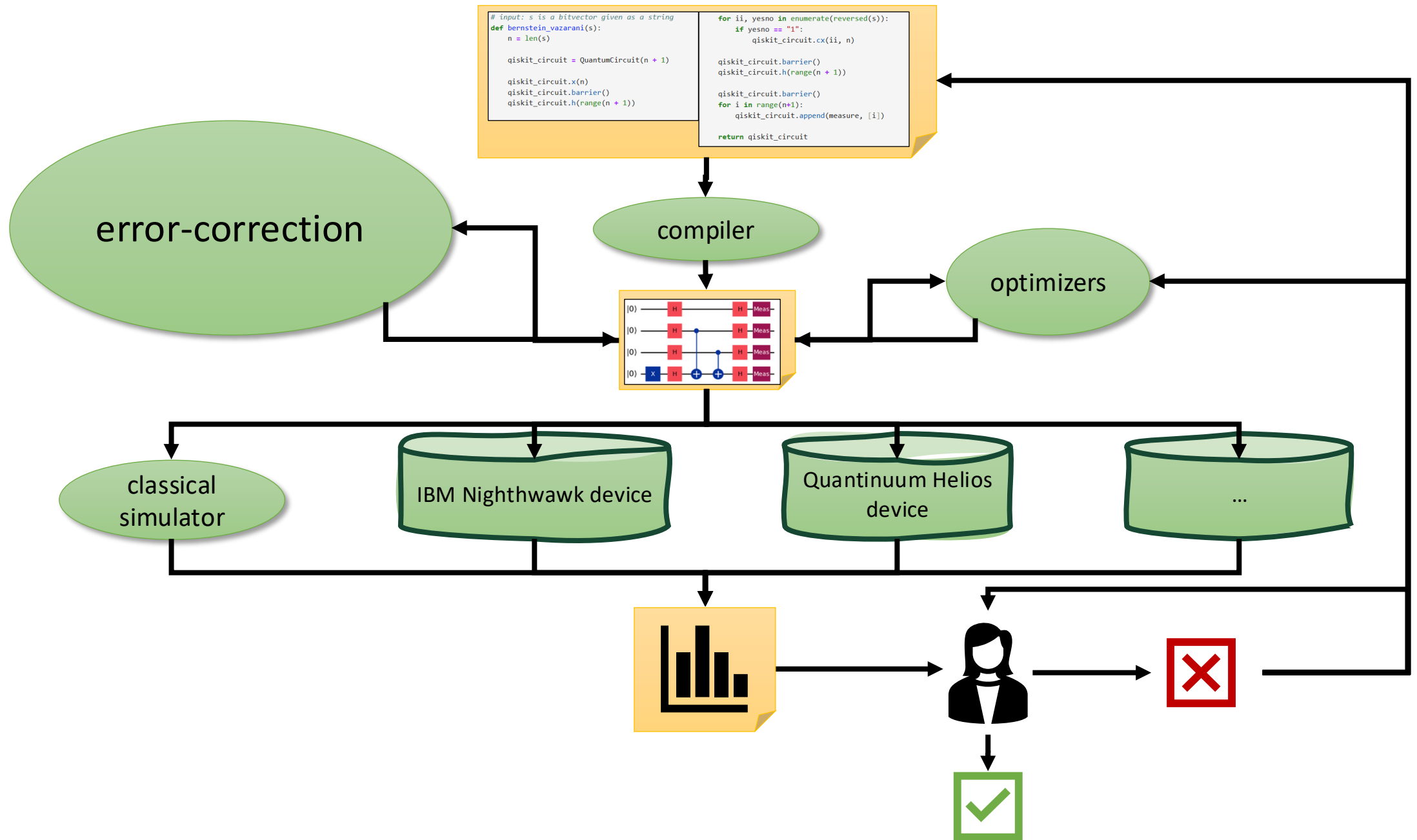
<sup>1</sup> University of Pennsylvania, Philadelphia, USA  
{keanchen, stevez, gushuli}@seas.upenn.edu

<sup>2</sup> University of Edinburgh, Edinburgh, UK

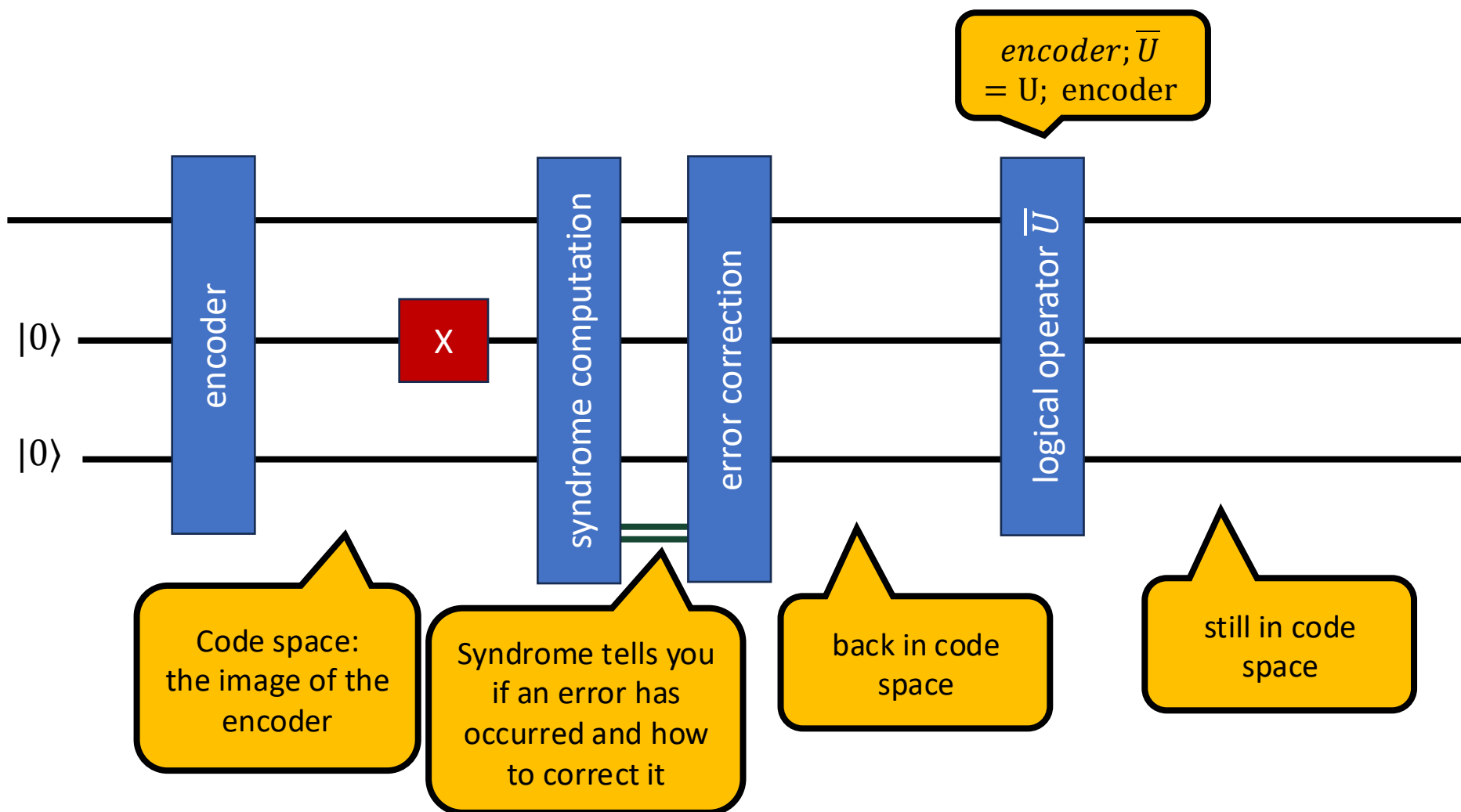
<sup>3</sup> Intel Corporation, Hillsboro, USA

<sup>4</sup> Intel Corporation, Santa Clara, USA





# Quantum error-correcting codes (QECC)



# Verifying QECCs

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## *QECV*: A VERIFICATION FRAMEWORK FOR QUANTUM ERROR CORRECTION CODES

Authors: Anbang Wu, Gushu Li, Hezi Zhang, Gian Giacomo Guerreschi, Yuan Xie, Yufei Ding

---

## Symbolic Execution for Quantum Error Correction Programs

Authors:  [Wang Fang](#),  [Mingsheng Ying](#) | [Authors Info & Claims](#)

[Proceedings of the ACM on Programming Languages, Volume 8, Issue PLDI](#) • Article No.: 189, Pages 1040 - 1065  
<https://doi.org/10.1145/3656419>

## Efficient Formal Verification of Quantum Error Correcting Programs

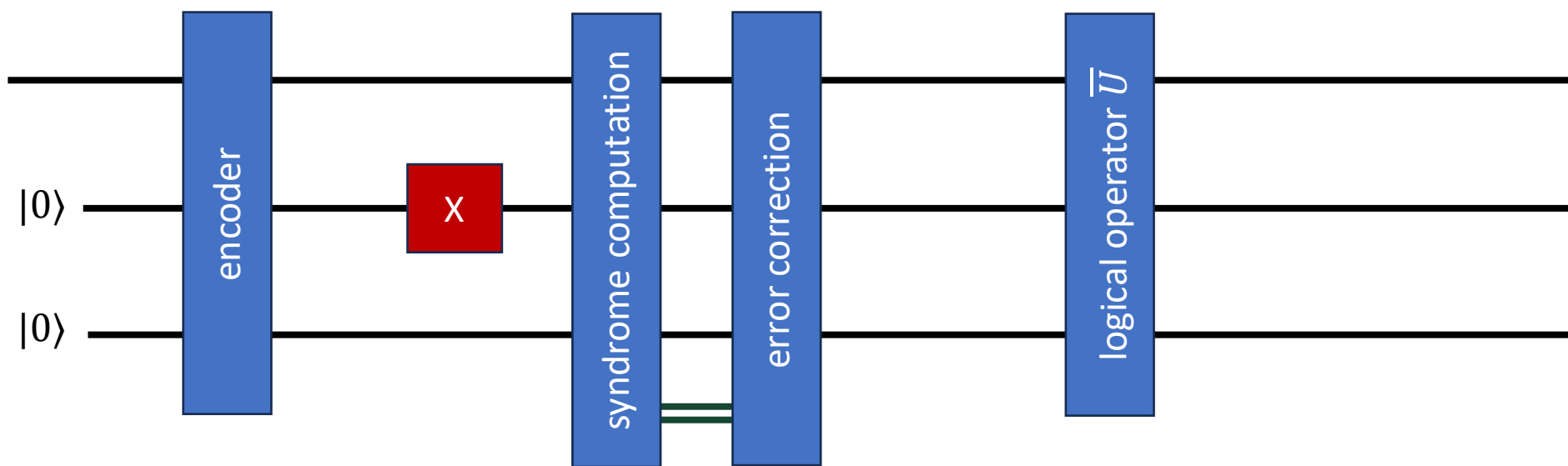
Authors:  [Qifan Huang](#),  [Li Zhou](#),  [Wang Fang](#),  [Mengyu Zhao](#),  [Mingsheng Ying](#) | [Authors Info & Claims](#)

[Proceedings of the ACM on Programming Languages, Volume 9, Issue PLDI](#) • Article No.: 190, Pages 1068 - 1093  
<https://doi.org/10.1145/3729293>

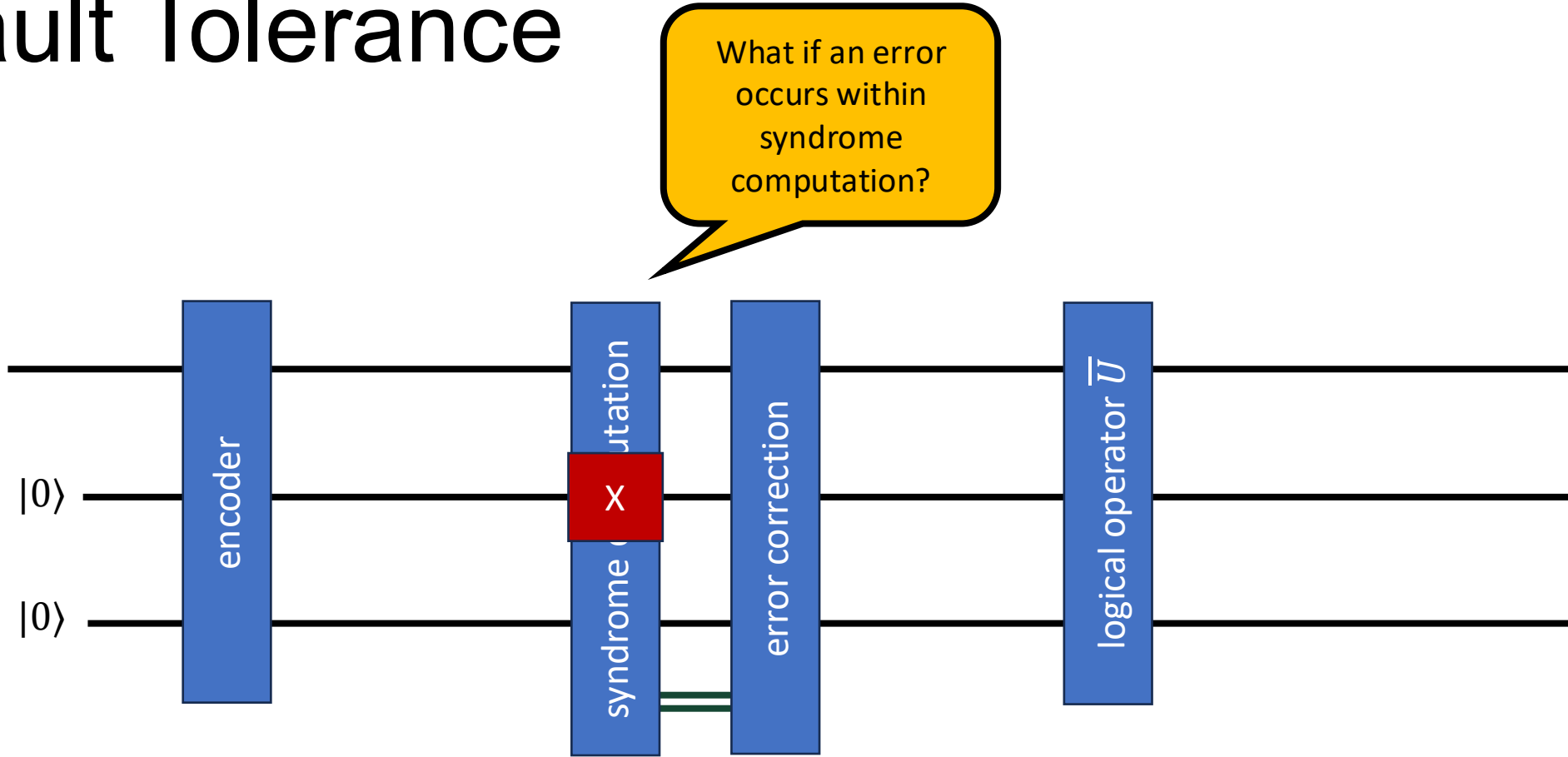
# Stabilizer Formalism

Qubits	Semantics
pure state $ \varphi\rangle \in \mathbb{C}^{2^n}$	$U \in \mathbb{C}^{2^n \times 2^n}$
mixed state <i>(probability distribution over pure states)</i> $\rho = \sum_i p_i  \varphi_i\rangle\langle\varphi_i  \in \mathbb{C}^{2^n \times 2^n}$	completely positive maps (CPMs) $\rho \mapsto \rho'$ i.e. $\text{tr}(\rho) \geq 0 \rightarrow \text{tr}(\rho') \geq 0$
stabilizer states <i>(set of Pauli operators that stabilize the state)</i> $S \subseteq \mathcal{P}_n \cong \mathbb{Z}_4 \times \mathbb{Z}_2^{2n}$ $\Rightarrow$ $\{  \varphi\rangle \mid \forall P \in S, P \varphi\rangle =  \varphi\rangle \}$	Clifford operators as Pauli tableaux: $T \in (\mathcal{P}_n)^{2n}$ <div style="display: flex; justify-content: space-around; margin-top: 20px;"> <div data-bbox="1480 1119 2033 1339" style="border: 2px solid black; border-radius: 50%; padding: 10px; background-color: yellow;">                         bitvector representation                     </div> <div data-bbox="1931 929 2548 1219" style="border: 2px solid black; border-radius: 50%; padding: 10px; background-color: yellow;">                         efficiently simulatable on a classical computer                     </div> </div>

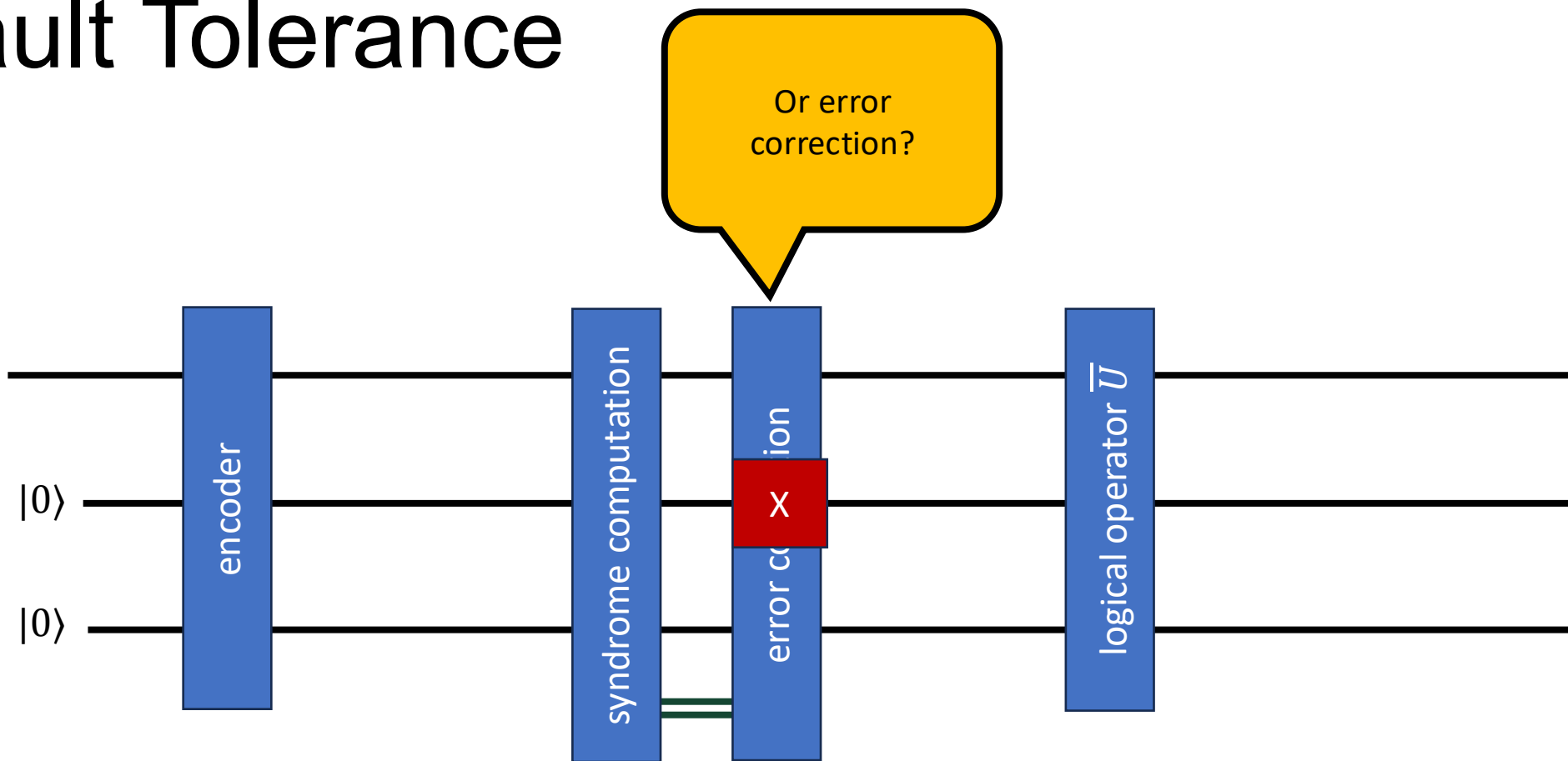
# Fault Tolerance



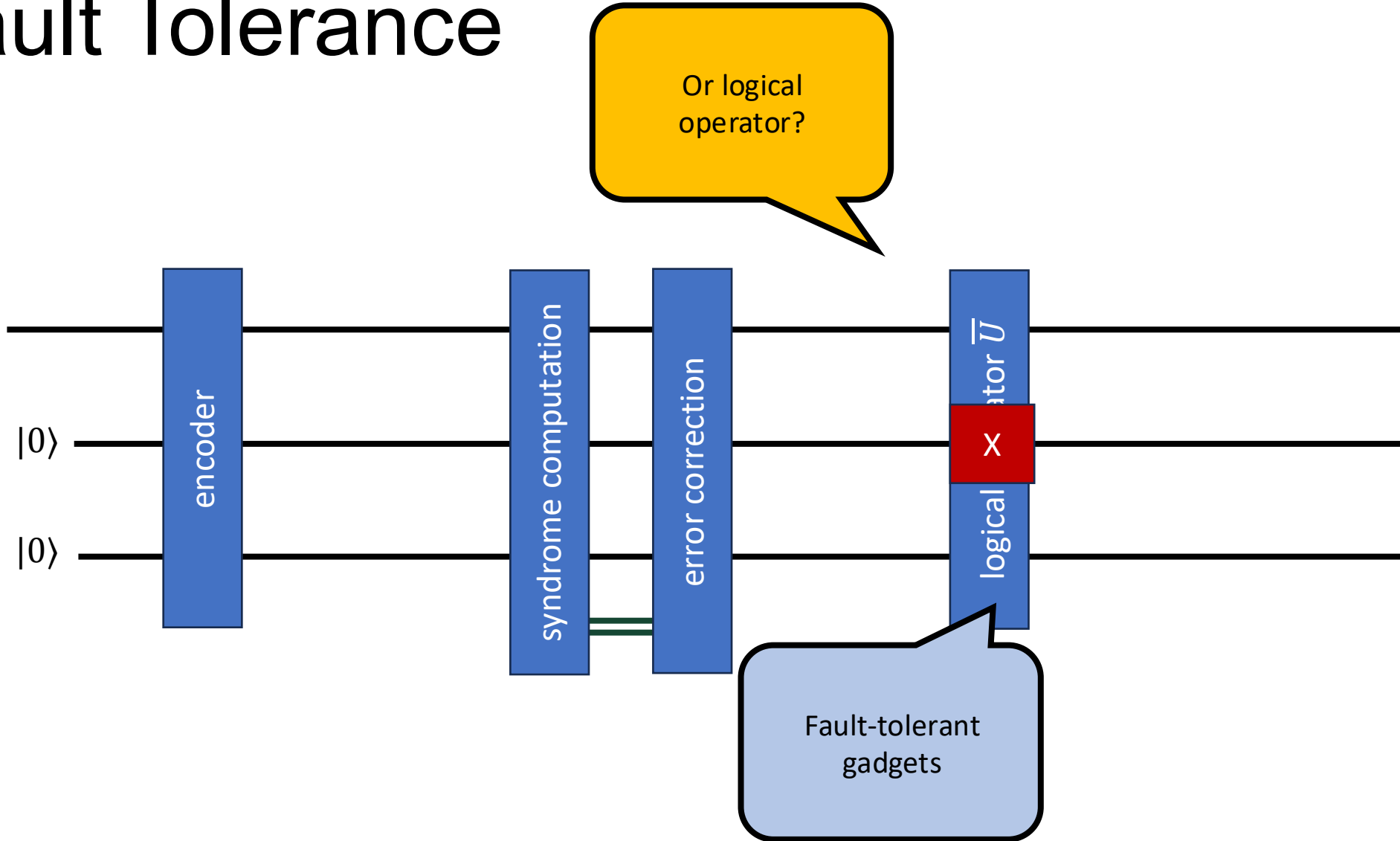
# Fault Tolerance



# Fault Tolerance



# Fault Tolerance



# Verifying fault-tolerance of gadgets

- Formalize fault-tolerance
- Symbolic execution with added quantum faults
- Special symbolic execution rules for repeat-until-success
- Plus: magic state distillation
  - Decompose rules into parts to be checked with above techniques

Check for updates

## Verifying Fault-Tolerance of Quantum Error Correction Codes

Kean Chen<sup>1</sup>(✉) , Yuhao Liu<sup>1</sup> , Wang Fang<sup>2</sup> , Jennifer Paykin<sup>3</sup>,  
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<sup>1</sup> University of Pennsylvania, Philadelphia, USA  
{keanchen, stevez, gushuli}@seas.upenn.edu

<sup>2</sup> University of Edinburgh, Edinburgh, UK

<sup>3</sup> Intel Corporation, Hillsboro, USA

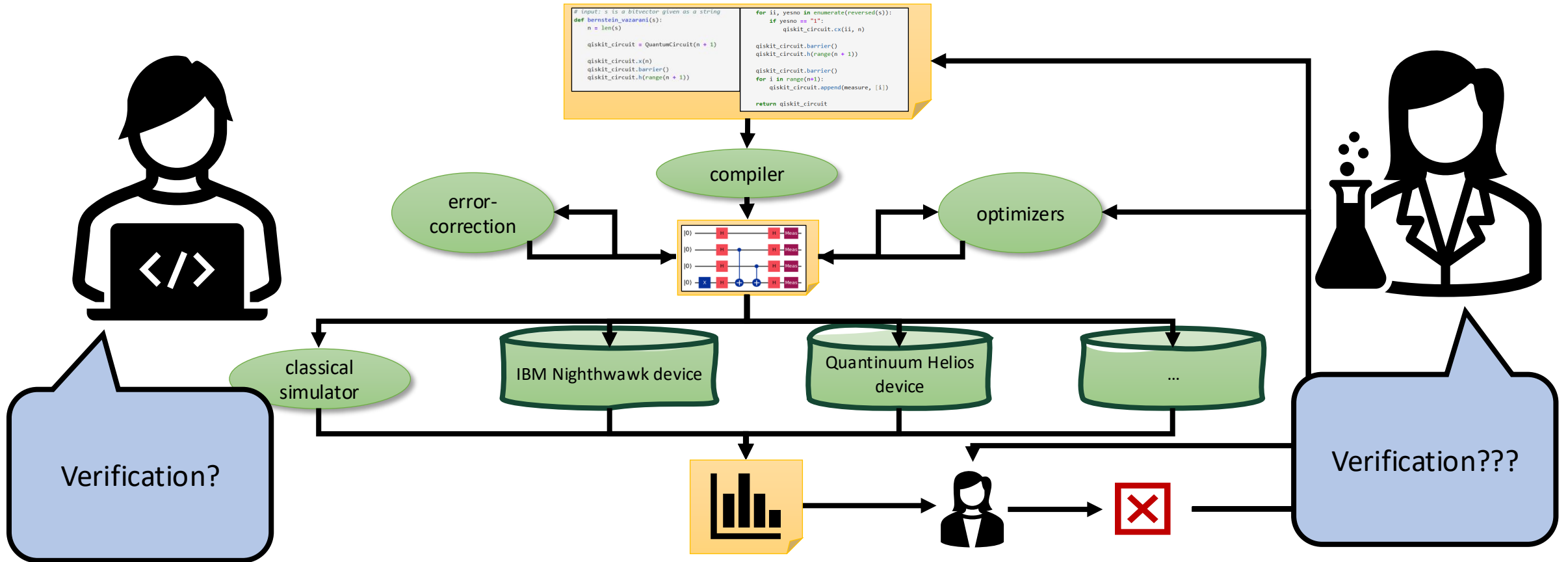
<sup>4</sup> Intel Corporation, Santa Clara, USA



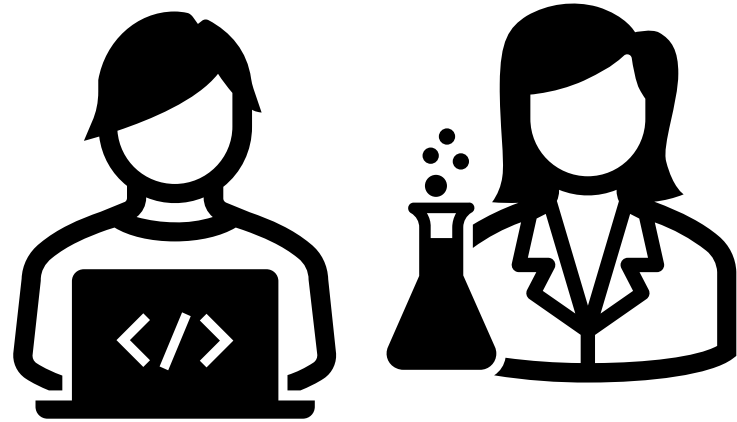
# Open problems

1. Structures to efficiently represent quantum semantics
2. Verifying more than circuits: quantum programs
3. Automation/expressive power tradeoff
4. How to intuitively express quantum specifications?
5. Verification in the presence of errors
  
6. Encourage adoption from all

# Open problem #6: Adoption and collaboration



# Open problem #6: Adoption and collaboration



Verification!

# Specification, Semantics, and Verification of Quantum Programs

Questions?

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Hiring  
PhD  
students!

Certified Programs and Proofs (CPP)  
Jan 13, 2026

